



# NSCET E-LEARNING PRESENTATION

LISTEN ... LEARN... LEAD...





# **ELECTRONICS AND COMMUNICATION ENGINEERING**

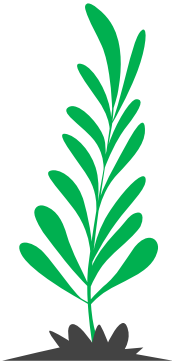
**III YEAR / VI th SEMESTER**

**EC8095– VLSI DESIGN**

**Ms.Vivitha,M.E**

**Assistant Professor**

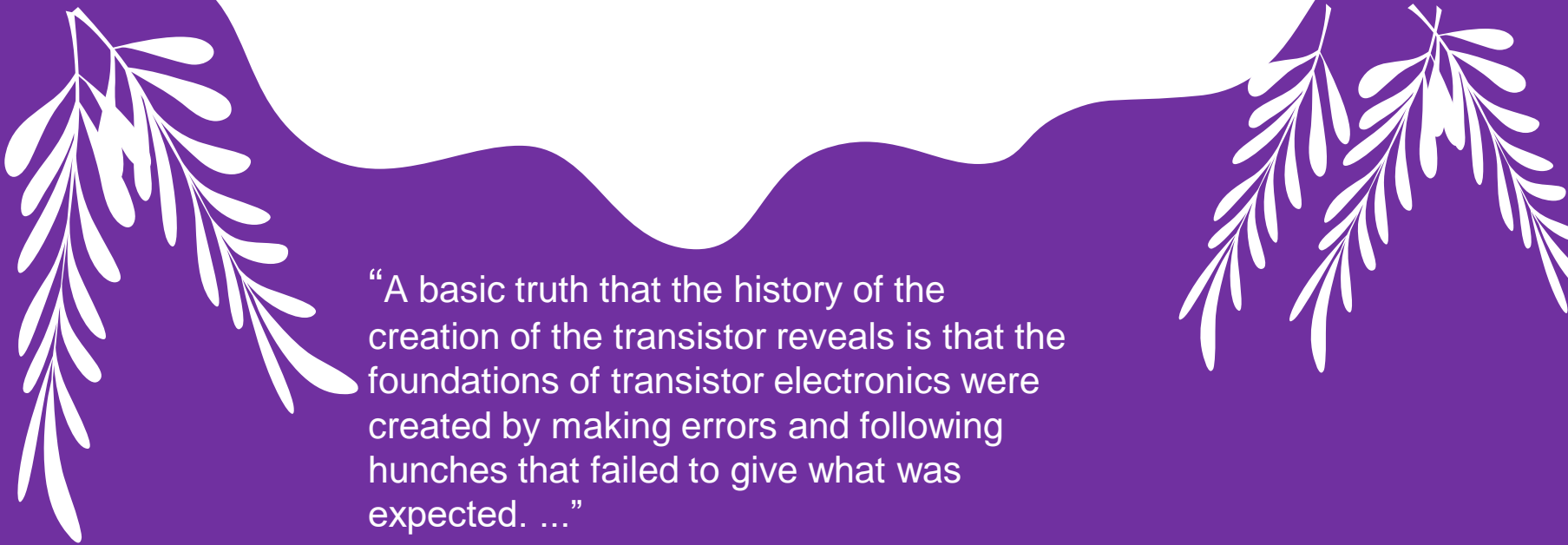
**Nadar Saraswathi College of & Technology,  
Vadapudupatti, Annanji (po), Theni – 625531.**





**UNIT II COMBINATIONAL MOS LOGIC CIRCUITS– LECTURE 02**





“A basic truth that the history of the creation of the transistor reveals is that the foundations of transistor electronics were created by making errors and following hunches that failed to give what was expected. ...”

**-William Shockley**



# 01 INTRODUCTION

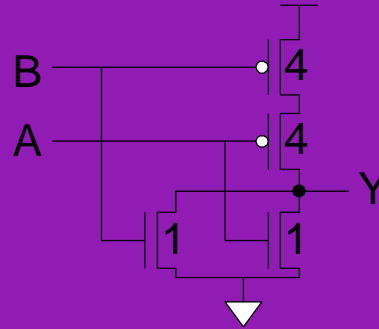
# 02 CIRCUIT FAMILIES

# 03 POWER

# 01. INTRODUCTION

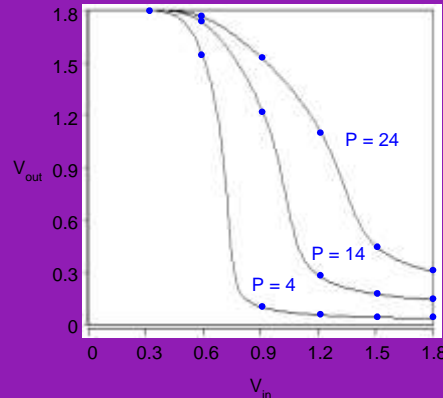
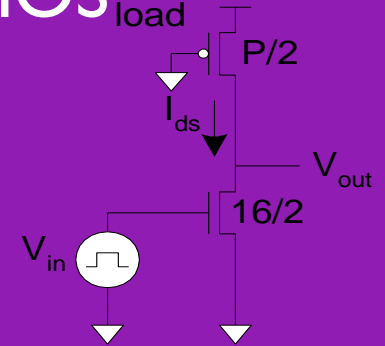
# COMBINATIONAL CIRCUIT

- What makes a circuit fast?
  - $I = C \, dV/dt \rightarrow t_{pd} \propto (C/I) \Delta V$
  - low capacitance
  - high current
  - small swing
- Logical effort is proportional to  $C/I$
- pMOS are the enemy!
  - High capacitance for a given current
- Can we take the pMOS capacitance off the input?
- Various circuit families try to do this...



# Pseudo-nMOS

- In the old days, nMOS processes had no pMOS
  - Instead, use pull-up transistor that is always ON
- In CMOS, use a pMOS that is always ON
  - *Ratio* issue
  - Make pMOS about  $1/4$  effective strength of pulldown network



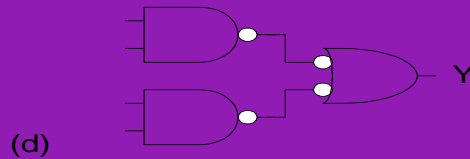
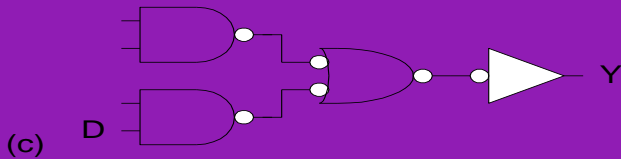
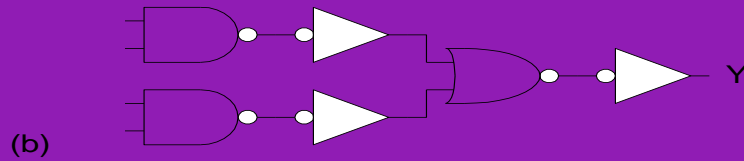
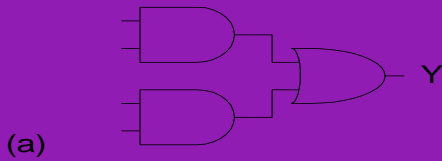


# STATIC CMOS

- Bubble Pushing
- Compound Gates
- Logical Effort Example
- Input Ordering
- Asymmetric Gates
- Skewed Gates
- Best P/N ratio

# Bubble Pushing

- Start with network of AND / OR gates
- Convert to NAND / NOR + inverters
- Push bubbles around to simplify logic
  - Remember DeMorgan's Law

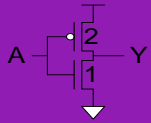


# Compound Gates

- Logical Effort of compound gates

unit inverter

$$Y = \overline{A}$$

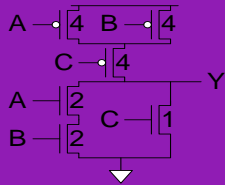


$$g_A = 3/3$$

$$p = 3/3$$

AOI21

$$Y = \overline{A \square B + C}$$



$$g_A = 6/3$$

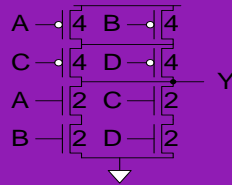
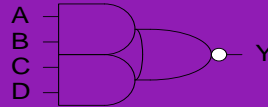
$$g_B = 6/3$$

$$g_C = 5/3$$

$$p = 7/3$$

AOI22

$$Y = \overline{A \square B + C \square D}$$



$$g_A = 6/3$$

$$g_B = 6/3$$

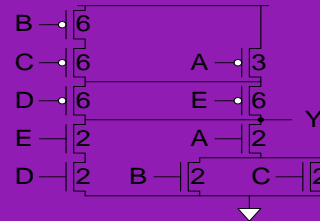
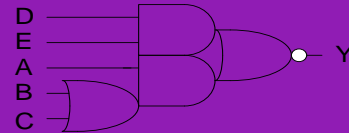
$$g_C = 6/3$$

$$g_D = 6/3$$

$$p = 12/3$$

Complex AOI

$$Y = \overline{A \square (B + C) + D \square E}$$



$$g_A = 5/3$$

$$g_B = 8/3$$

$$g_C = 8/3$$

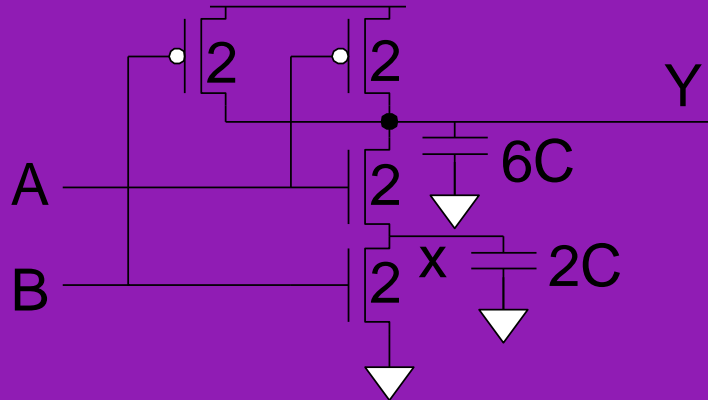
$$g_D = 8/3$$

$$g_E = 8/3$$

$$p = 16/3$$

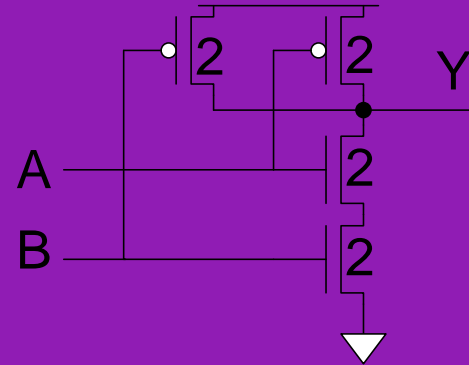
# Input Order

- Our parasitic delay model was too simple
  - Calculate parasitic delay for Y falling
    - If A arrives latest?  $2\tau$
    - If B arrives latest?  $2.33\tau$



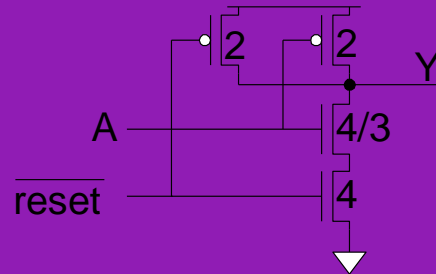
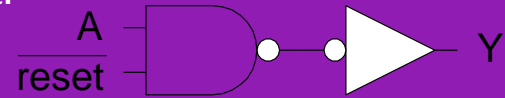
# Inner & Outer Inputs

- *Inner* input is closest to output (A)
- *Outer* input is closest to rail (B)
- If input arrival time is known
  - Connect latest input to inner terminal



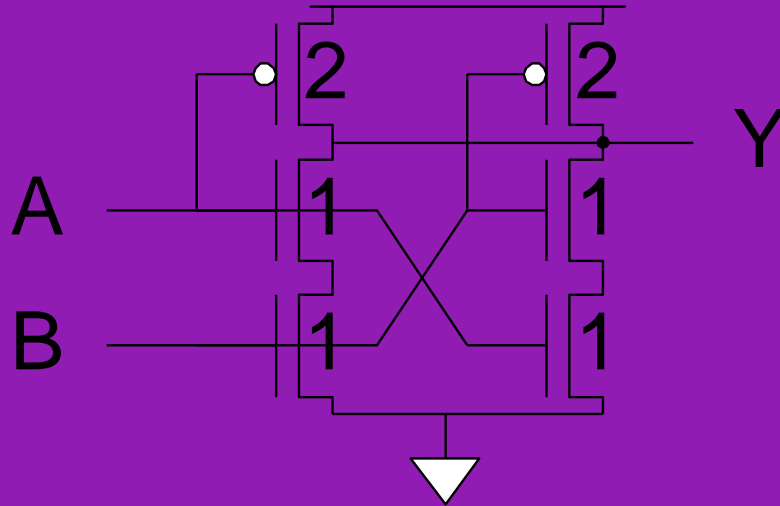
# Asymmetric Gates

- Asymmetric gates favor one input over another
- Ex: suppose input A of a NAND gate is most critical
  - Use smaller transistor on A (less capacitance)
  - Boost size of noncritical input
  - So total resistance is same
- $g_A = 10/9$
- $g_B = 2$
- $g_{\text{total}} = g_A + g_B = 28/9$
- Asymmetric gate approaches  $g = 1$  on critical input
- But total logical effort goes up



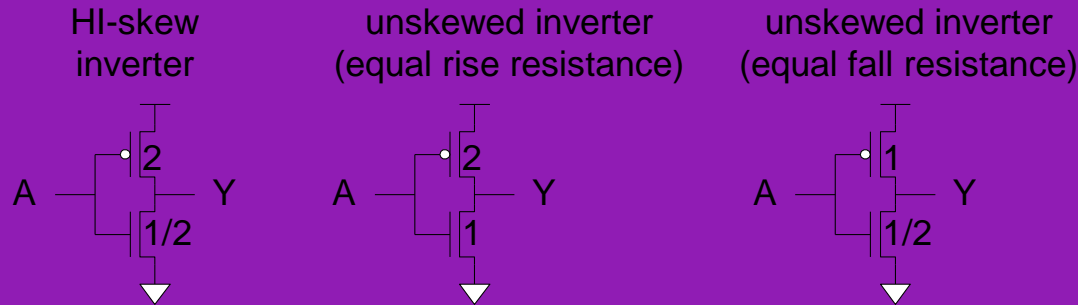
# Symmetric Gates

- Inputs can be made perfectly symmetric



# Skewed Gates

- Skewed gates favor one edge over another
- Ex: suppose rising output of inverter is most critical
  - Downsize noncritical nMOS transistor



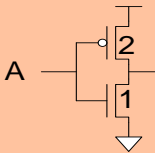
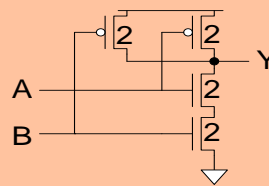
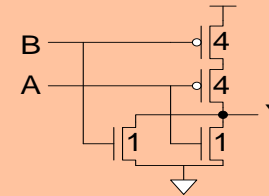
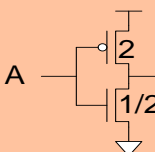
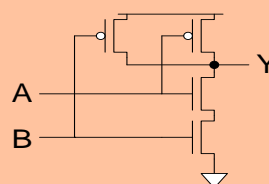
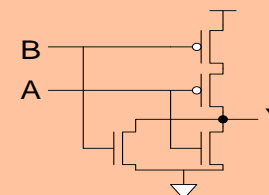
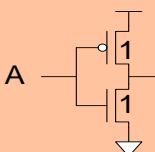
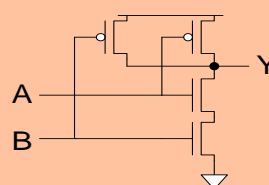
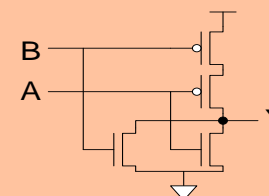
- Calculate logical effort by comparing to unskewed inverter with same effective resistance on that edge.
  - $g_u = 2.5 / 3 = 5/6$
  - $g_d = 2.5 / 1.5 = 5/3$



# HI- and LO-Skew

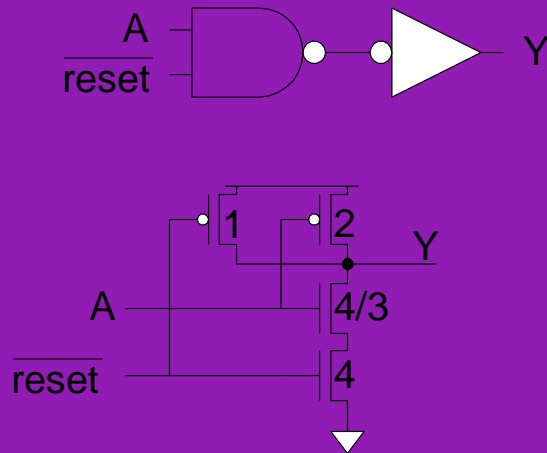
- Def: Logical effort of a skewed gate for a particular transition is the ratio of the input capacitance of that gate to the input capacitance of an unskewed inverter delivering the same output current for the same transition.
- Skewed gates reduce size of noncritical transistors
  - HI-skew gates favor rising output (small nMOS)
  - LO-skew gates favor falling output (small pMOS)
- Logical effort is smaller for favored direction
- But larger for the other direction

# Catalog of Skewed Gates

|          | Inverter  | NAND2  | NOR2  |
|----------|---|--|---|
| unskewed |  $g_u = 2$ $g_d = 1$ $g_{avg} = 1$     |  $g_u = 2$ $g_d = 2$ $g_{avg} = 4/3$ |  $g_u = 4$ $g_d = 1$ $g_{avg} = 5/3$ |
| HI-skew  |  $g_u = 2$ $g_d = 1/2$ $g_{avg} = 5/4$ |  $g_u = 2$ $g_d = 1$ $g_{avg} = 3/2$ |  $g_u = 4$ $g_d = 1$ $g_{avg} = 5/3$ |
| LO-skew  |  $g_u = 1$ $g_d = 1$ $g_{avg} = 1$     |  $g_u = 1$ $g_d = 1$ $g_{avg} = 1$   |  $g_u = 1$ $g_d = 1$ $g_{avg} = 1$   |

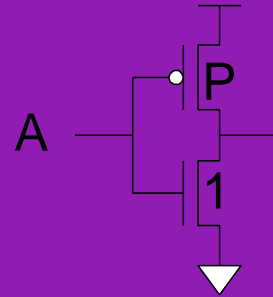
# Asymmetric Skew

- Combine asymmetric and skewed gates
  - Downsize noncritical transistor on unimportant input
  - Reduces parasitic delay for critical input



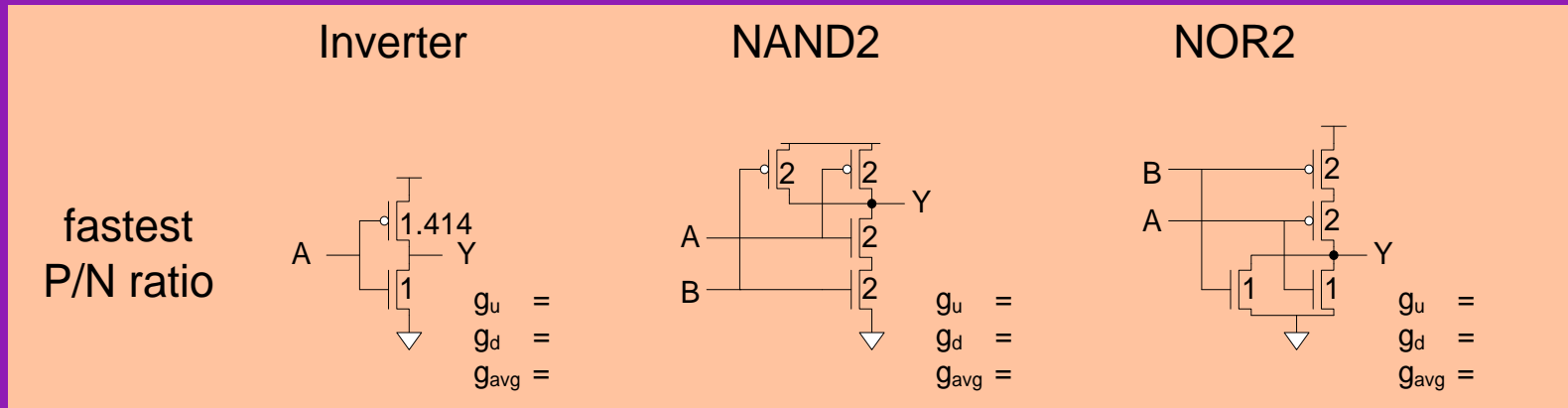
# Best P/N Ratio

- We have selected P/N ratio for unit rise and fall resistance ( $\mu = 2-3$  for an inverter).
- Alternative: choose ratio for least average delay
- Ex: inverter
  - Delay driving identical inverter
  - $t_{pdf} = (P+1)$
  - $t_{pdr} = (P+1)(\mu/P)$
  - $t_{pd} = (P+1)(1+\mu/P)/2 = (P + 1 + \mu + \mu/P)/2$
  - $dt_{pd} / dP = (1 - \mu/P^2)/2 = 0$
  - Least delay for P



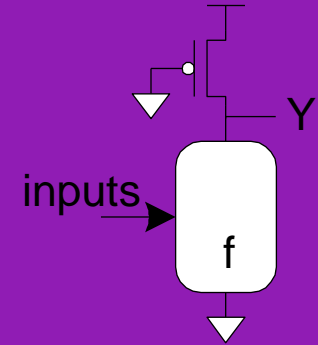
# P/N Ratios

- In general, best P/N ratio is sqrt of equal delay ratio.
  - Only improves average delay slightly for inverters
  - But significantly decreases area and power

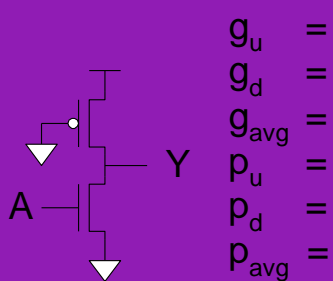


# Pseudo-nMOS Gates

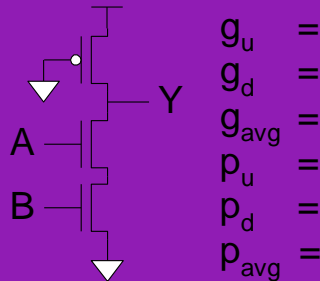
- Design for unit current on output to compare with unit inverter.
- pMOS fights nMOS



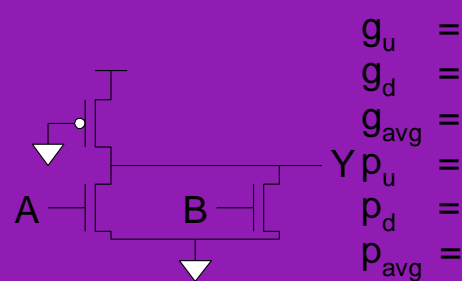
Inverter



NAND2



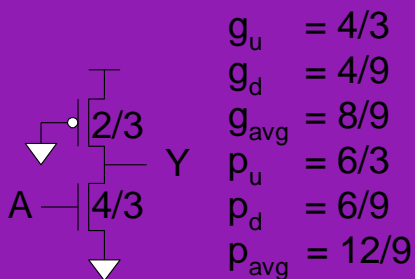
NOR2



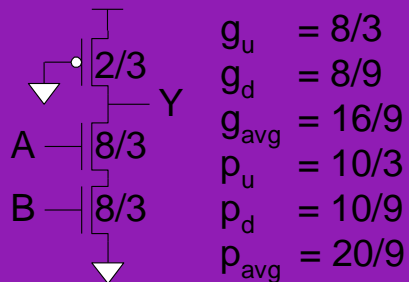
# Pseudo-nMOS Gates

- ❖ Design for unit current on output to compare with unit inverter.
- ❖ pMOS fights nMOS

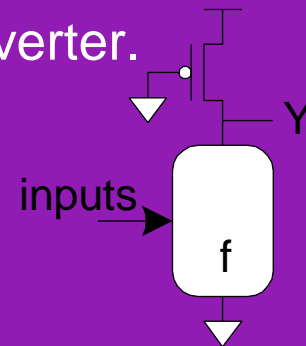
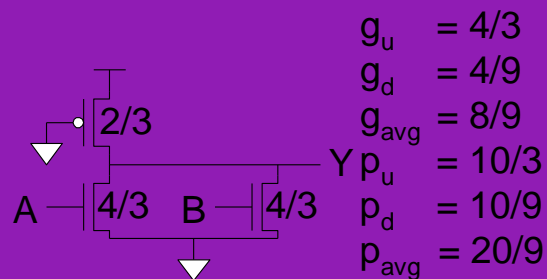
Inverter



NAND2

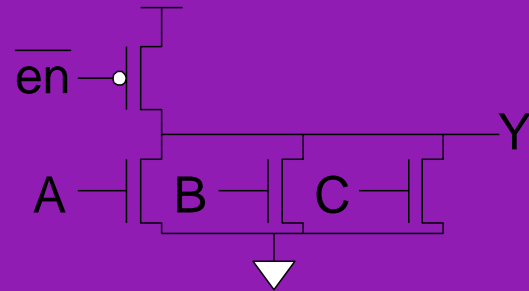


NOR2



# Pseudo-nMOS Power

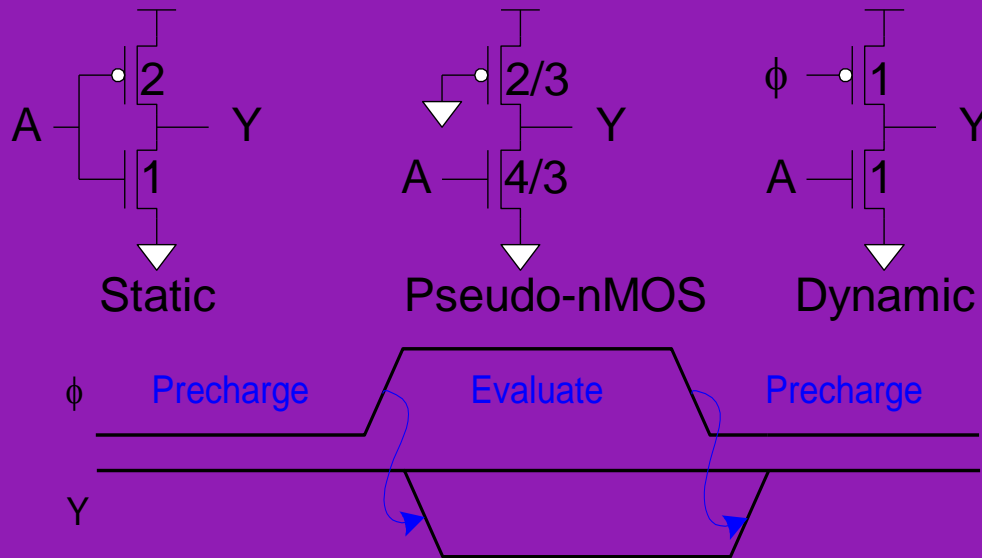
- Pseudo-nMOS draws power whenever  $Y = 0$ 
  - Called static power  $P = I_{DD}V_{DD}$
  - A few mA / gate \* 1M gates would be a problem
  - Explains why nMOS went extinct
- Use pseudo-nMOS sparingly for wide NORs
- Turn off pMOS when not in use





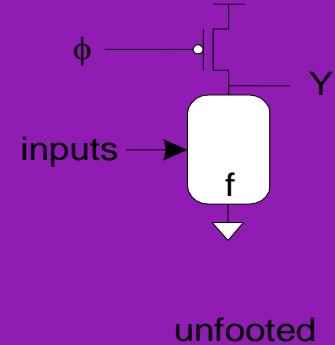
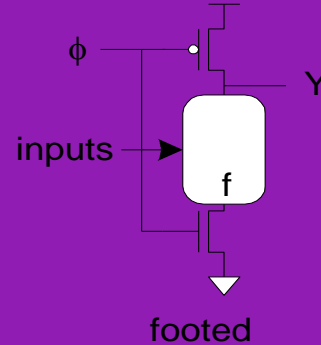
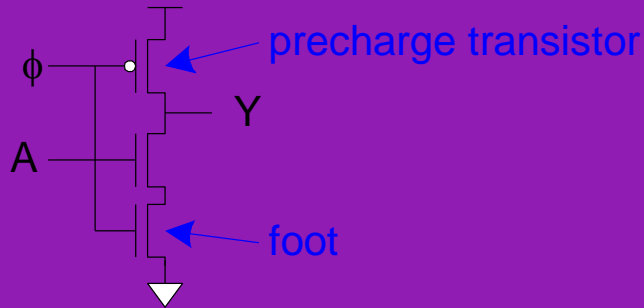
# Dynamic Logic

- *Dynamic gates* uses a clocked pMOS pullup
- Two modes: *precharge* and *evaluate*



# The Foot

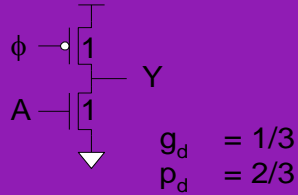
- What if pulldown network is ON during precharge?
- Use series evaluation transistor to prevent fight.



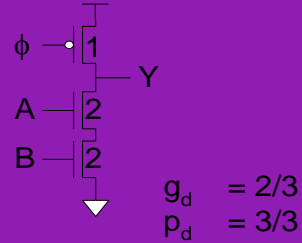
# Logical Effort

unfooted

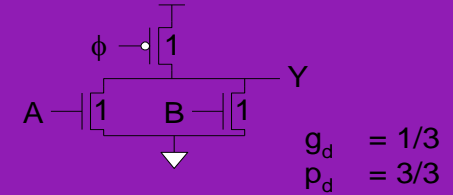
Inverter



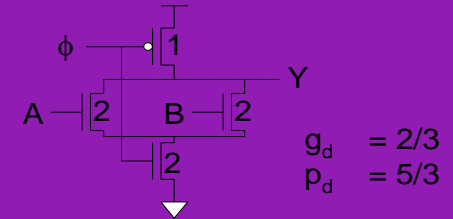
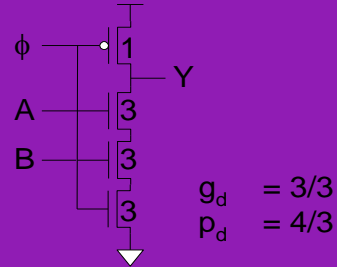
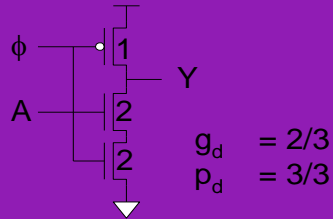
NAND2



NOR2

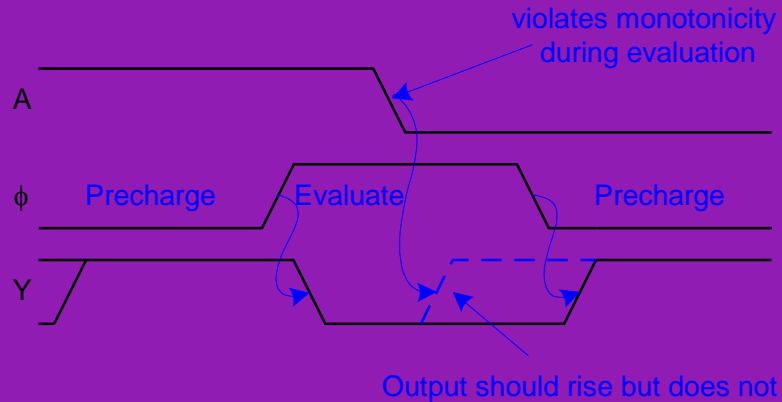
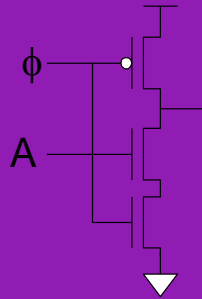


footed



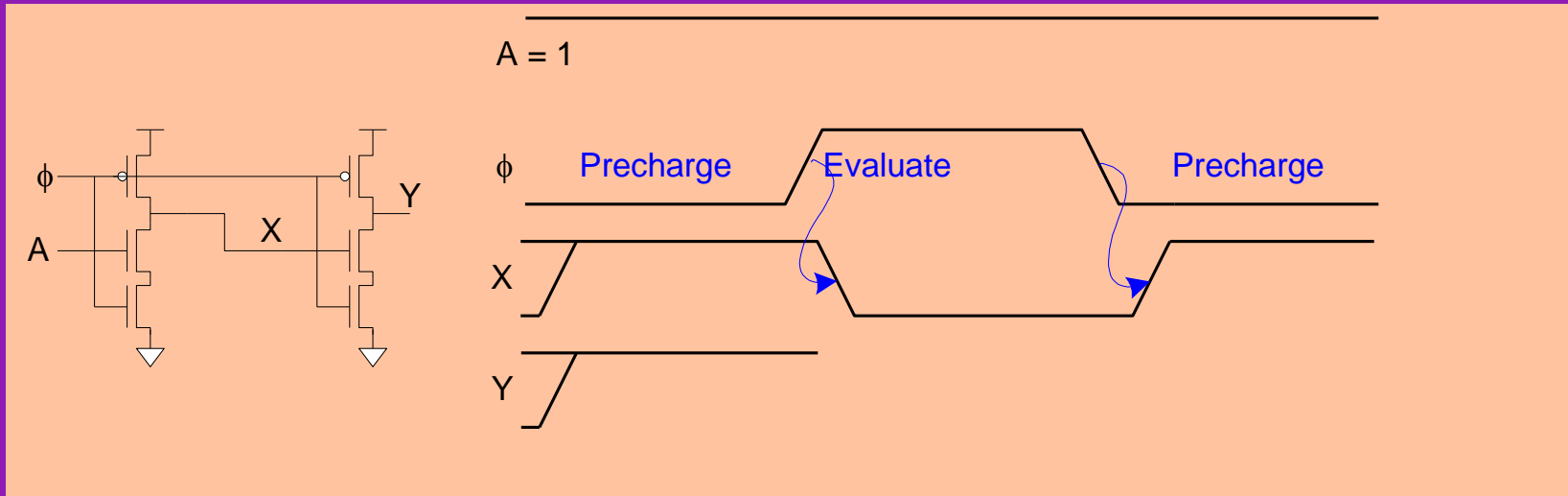
# Monotonicity

- Dynamic gates require *monotonically rising* inputs during evaluation
  - 0 -> 0
  - 0 -> 1
  - 1 -> 1
  - But not 1 -> 0



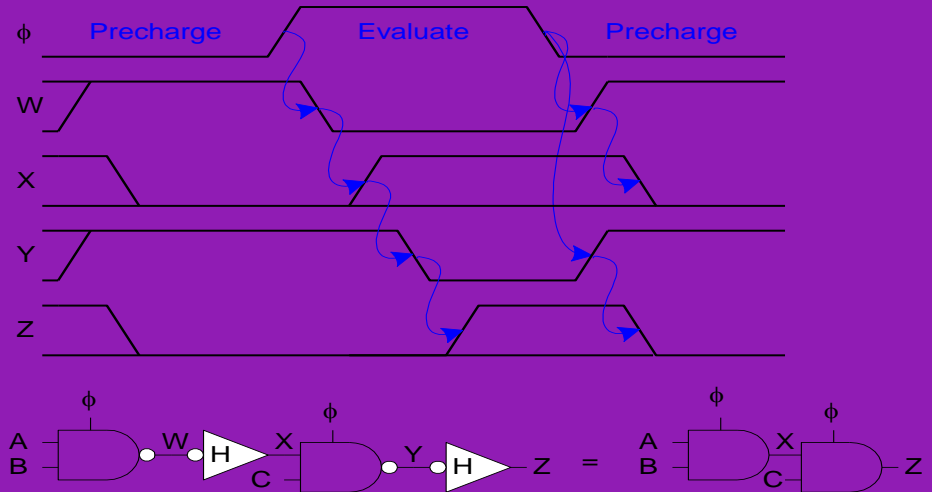
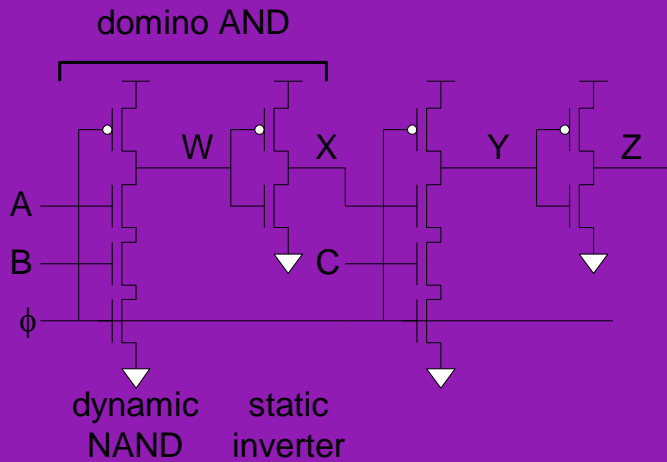
# Monotonicity Woes

- But dynamic gates produce monotonically falling outputs during evaluation
- Illegal for one dynamic gate to drive another!



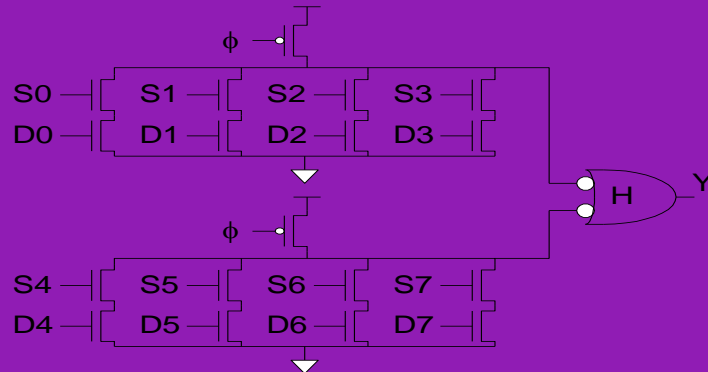
# Domino Gates

- Follow dynamic stage with inverting static gate
  - Dynamic / static pair is called domino gate
  - Produces monotonic outputs



# Domino Optimizations

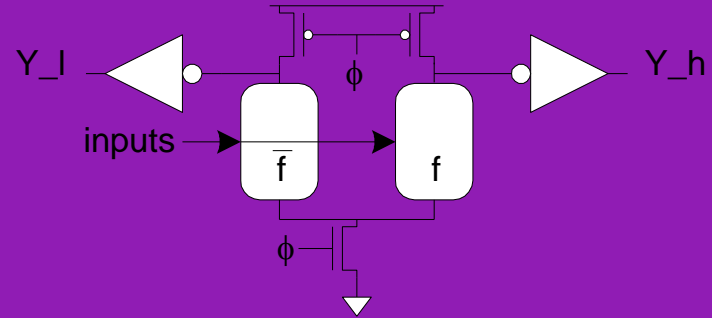
- Each domino gate triggers next one, like a string of dominos toppling over
- Gates evaluate sequentially but precharge in parallel
- Thus evaluation is more critical than precharge
- HI-skewed static stages can perform logic



# Dual-Rail Domino

- Domino only performs noninverting functions:
  - AND, OR but not NAND, NOR, or XOR
- Dual-rail domino solves this problem
  - Takes true and complementary inputs
  - Produces true and complementary outputs

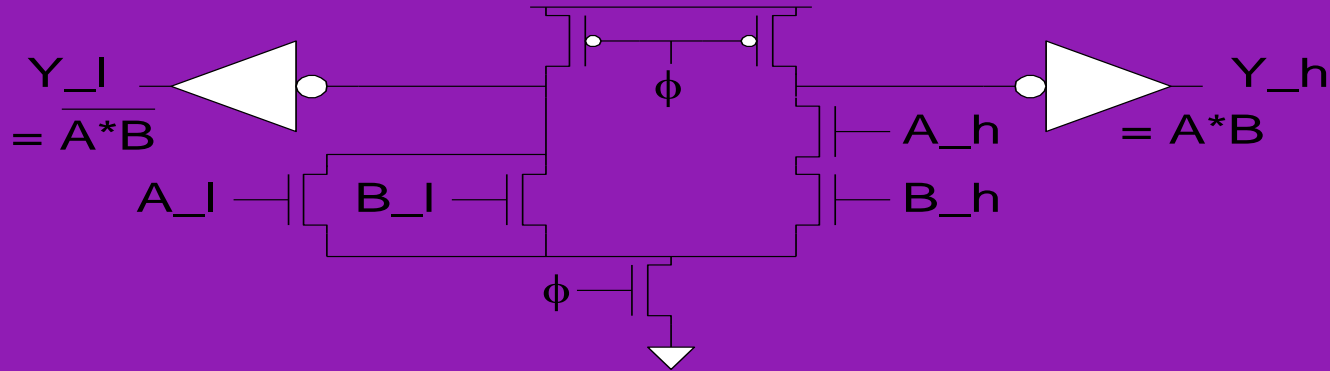
| sig_h | sig_l | Meaning    |
|-------|-------|------------|
| 0     | 0     | Precharged |
| 0     | 1     | '0'        |
| 1     | 0     | '1'        |
| 1     | 1     | invalid    |





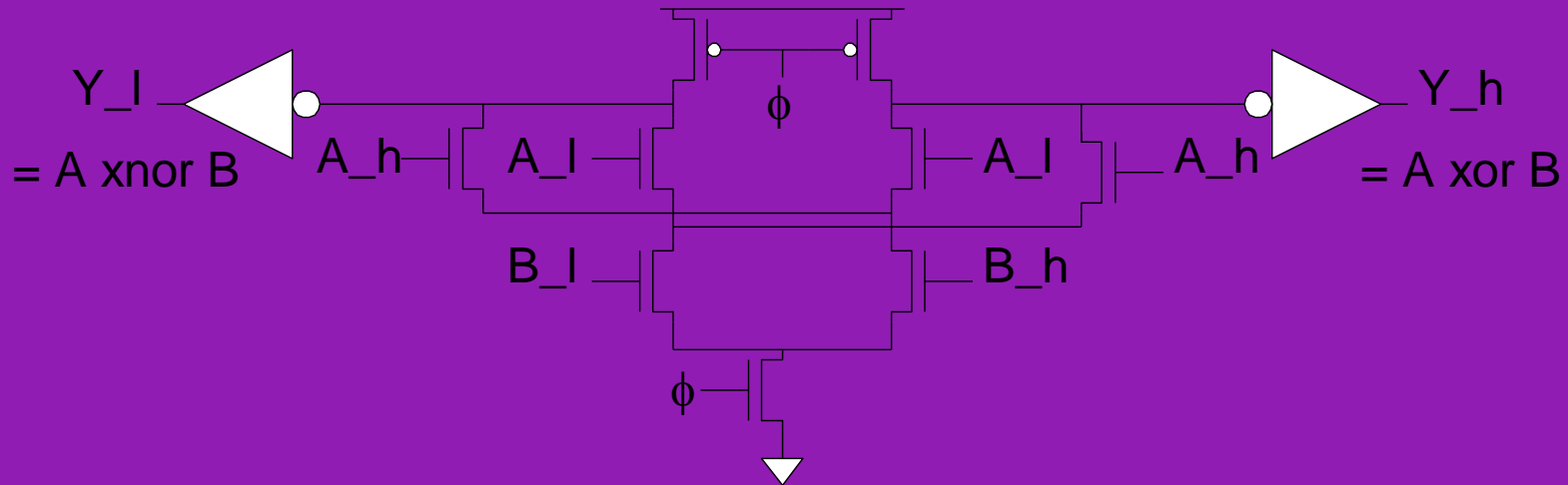
# Example: AND/NAND

- Given  $A\_h, A\_l, B\_h, B\_l$
- Compute  $Y\_h = AB, Y\_l = \overline{AB}$
- Pulldown networks are conduction complements



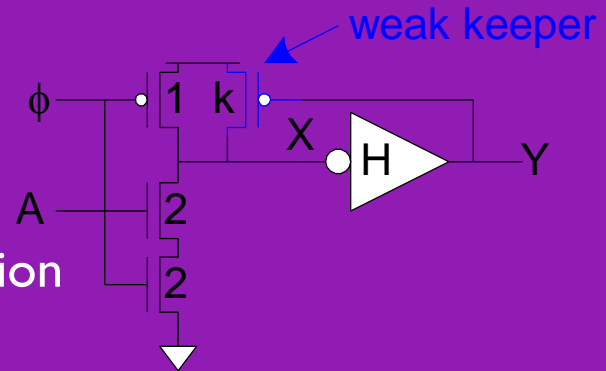
# Example: XOR/XNOR

- Sometimes possible to share transistors



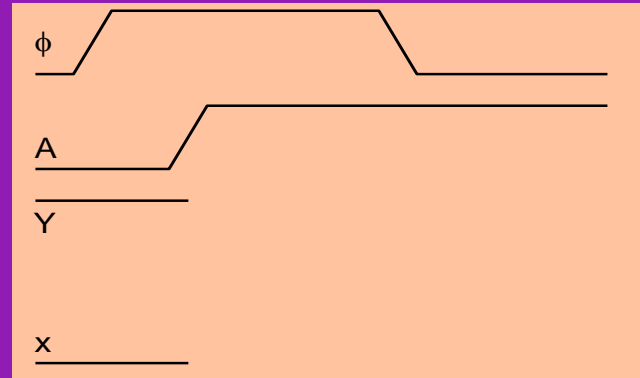
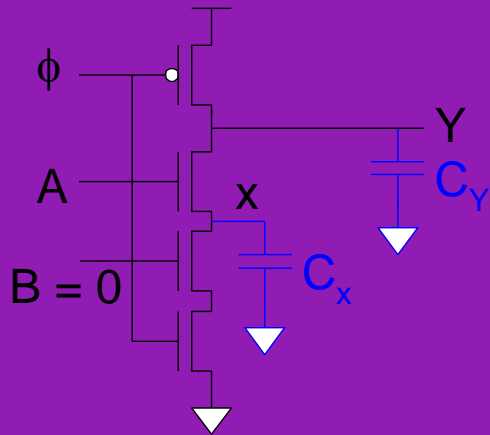
# Leakage

- Dynamic node floats high during evaluation
  - Transistors are leaky ( $I_{OFF} \neq 0$ )
  - Dynamic value will leak away over time
  - Formerly milliseconds, now nanoseconds
- Use keeper to hold dynamic node
  - Must be weak enough not to fight evaluation



# Charge Sharing

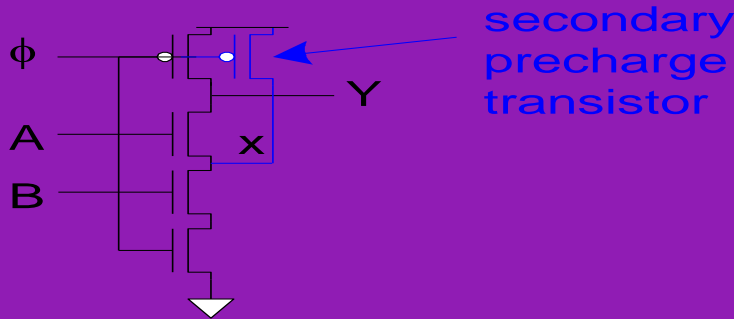
- Dynamic gates suffer from charge sharing



$$V_x = V_Y = \frac{C_Y}{C_x + C_Y} V_{DD}$$

# Secondary Precharge

- Solution: add secondary precharge transistors
  - Typically need to precharge every other node
- Big load capacitance  $C_Y$  helps as well



# Noise Sensitivity

- Dynamic gates are very sensitive to noise
  - Inputs:  $V_{IH} \approx V_{tn}$
  - Outputs: floating output susceptible noise
- Noise sources
  - Capacitive crosstalk
  - Charge sharing
  - Power supply noise
  - Feedthrough noise
  - And more!

# Power

- Domino gates have high activity factors
  - Output evaluates and precharges
    - If output probability = 0.5,  $\alpha = 0.5$ 
      - Output rises and falls on half the cycles
  - Clocked transistors have  $\alpha = 1$
- Leads to very high power consumption

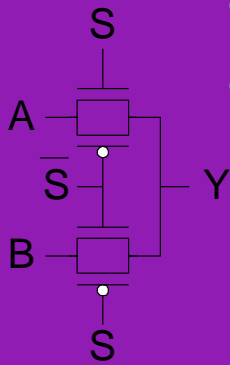
# Domino Summary

- Domino logic is attractive for high-speed circuits
  - 1.3 – 2x faster than static CMOS
  - But many challenges:
    - Monotonicity, leakage, charge sharing, noise
- Widely used in high-performance microprocessors in 1990s when speed was king
- Largely displaced by static CMOS now that power is the limiter
- Still used in memories for area efficiency

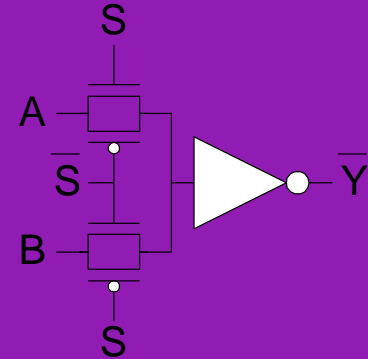


# Pass Transistor Circuits

- Use pass transistors like switches to do logic
- Inputs drive diffusion terminals as well as gates
- CMOS + Transmission Gates:

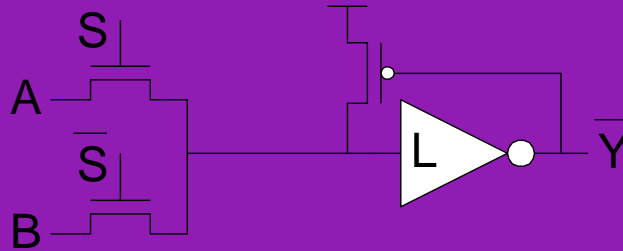


- 2-input multiplexer
- Gates should be restoring



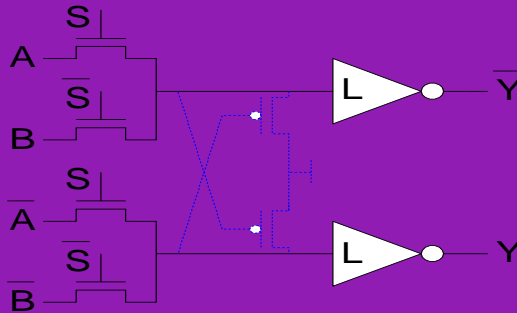
# LEAP

- **LEA**n integration with **P**ass transistors
- Get rid of pMOS transistors
  - Use weak pMOS feedback to pull fully high
  - Ratio constraint



# CPL

- **C**omplementary **P**ass-transistor **L**ogic
  - Dual-rail form of pass transistor logic
  - Avoids need for ratioed feedback
  - Optional cross-coupling for rail-to-rail swing



# Pass Transistor Summary

- Researchers investigated pass transistor logic for general purpose applications in the 1990's
  - Benefits over static CMOS were small or negative
  - No longer generally used
- However, pass transistors still have a niche in special circuits such as memories where they offer small size and the threshold drops can be managed

THANK YOU